

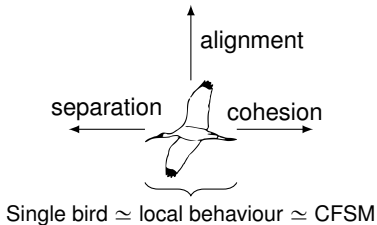
From Communicating Machines to Graphical Choreographies

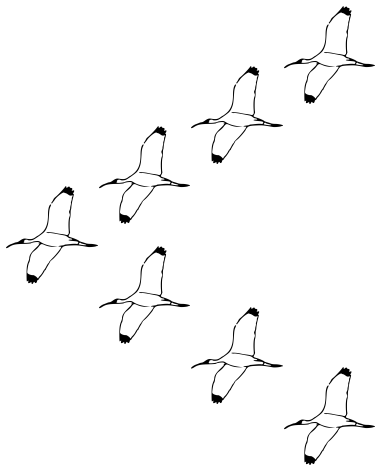
Julien Lange

joint work with Emilio Tuosto and Nobuko Yoshida

July 2015







Flock \simeq global behaviour \simeq choreography

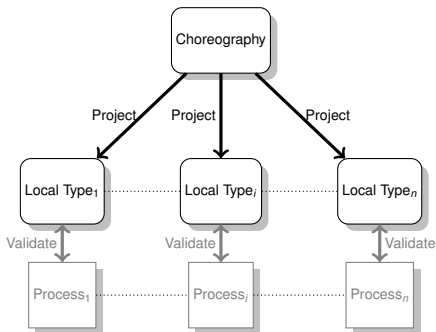


Introduction

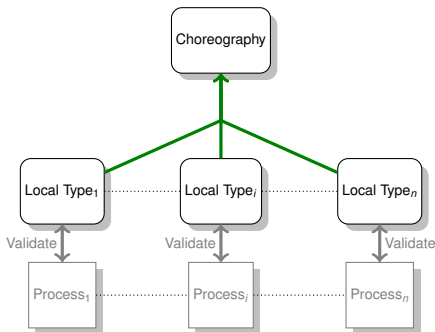
- ▶ Parts of distributed systems change/evolve, not always in a coordinated way,
- ▶ these changes are often *not* documented.
- ▶ Service oriented systems are sometimes composed dynamically,
- ▶ it is often unclear how complex the overall system has become.
- ▶ Cognizant's Zero Deviation Lifecycle Business Unit.

A *global* point of view of a distributed system is *essential* for top-level management.





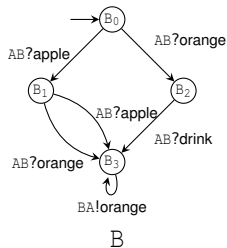
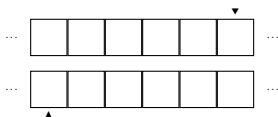
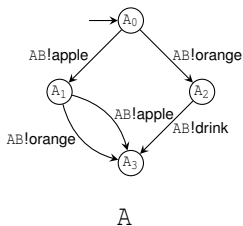
- ▶ Choreography-driven development, cf. Multipart Session Types top-down approach (POPL'08 & ESOP'12)
- ▶ Not applicable without *a priori knowledge* of a choreography



- ▶ Choreography-driven development, cf. Multiparty Session Types top-down approach (POPL'08 & ESOP'12)
- ▶ Not applicable without *a priori knowledge* of a choreography
- ▶ Our goal: from *Communicating Finite-State Machines* to *Global Graphs*



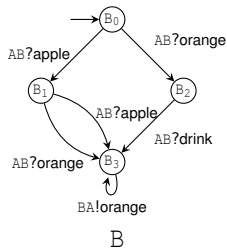
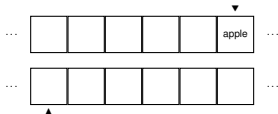
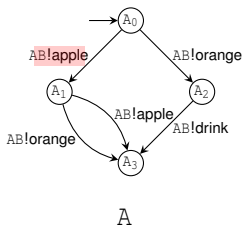
Background: CFSMs



“On Communicating Finite-State Machines”, Brand & Zafiropulo ('83)



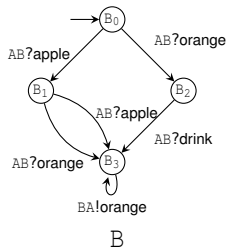
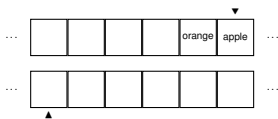
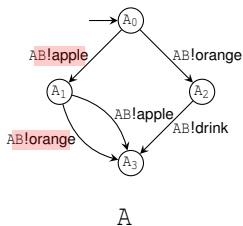
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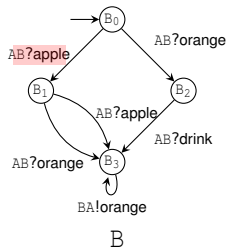
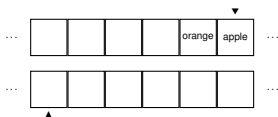
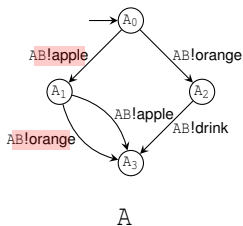
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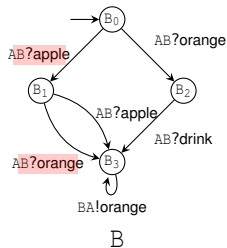
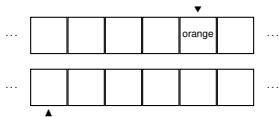
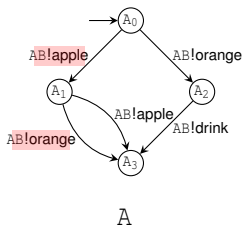
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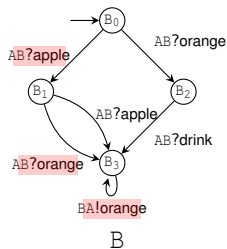
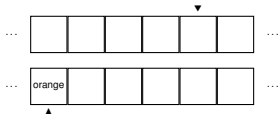
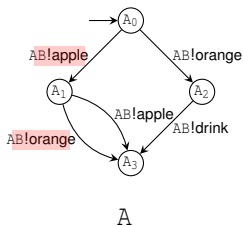
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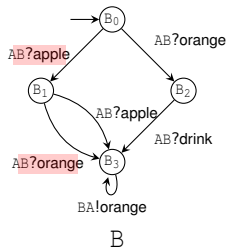
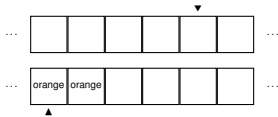
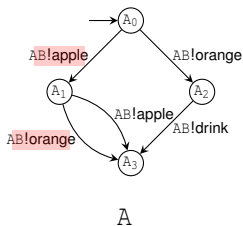
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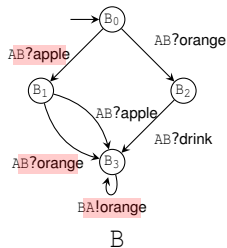
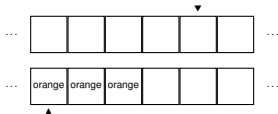
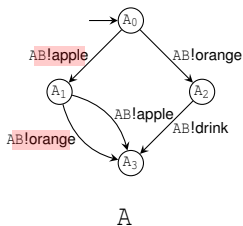
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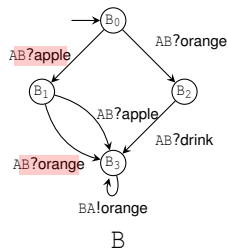
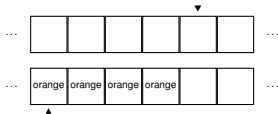
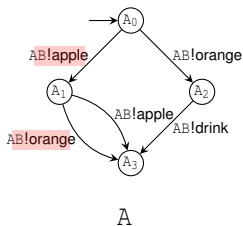
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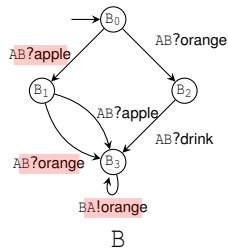
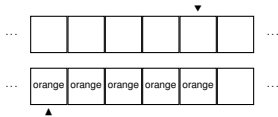
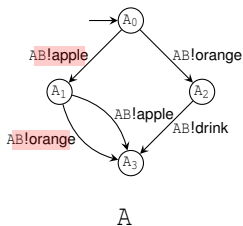
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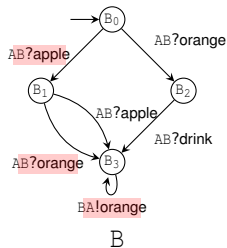
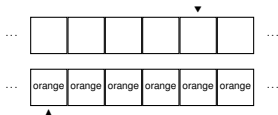
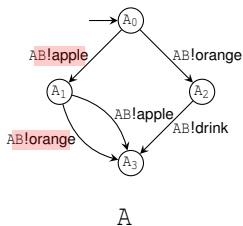
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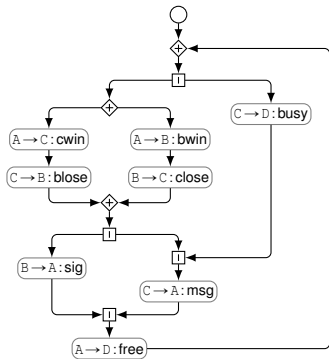
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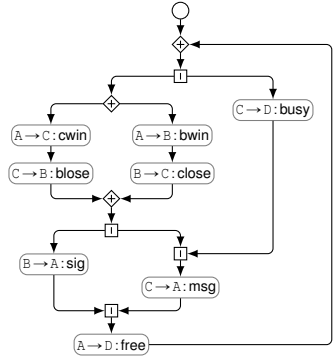
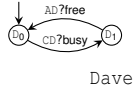
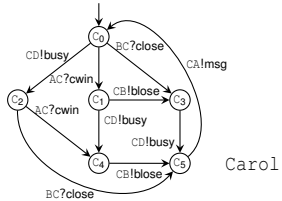
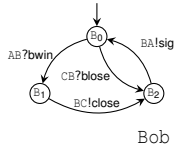
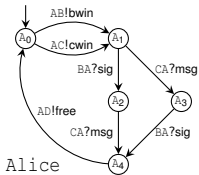


Global Graphs



Four Player Game

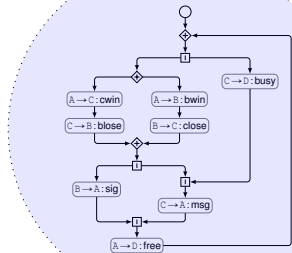
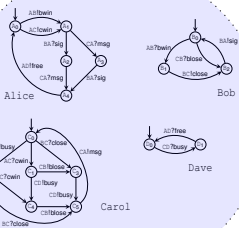
Global Graphs



Four Player Game



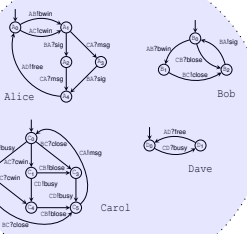
CFSMs



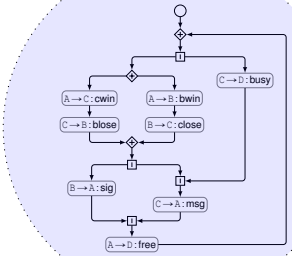
Deniélou & Yoshida – ESOP'12

CFSMs

This work



This work



Objectives

Two main objectives:

- ▶ **Sound Condition for Safety:** generalised multiparty compatibility

If $S = (M_1, \dots, M_k)$ is *compatible* then S is “safe”, i.e., every sent message is eventually received and no deadlock.

- ▶ **Construction of a Global Graph:**

If G is the global graph constructed from S , then

$$S = (M_1, \dots, M_k) \equiv (G \downarrow_1, \dots, G \downarrow_k)$$



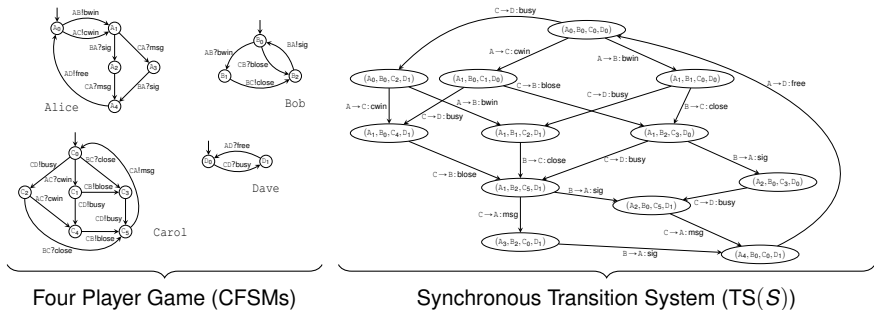
The Plan

1. Build $TS(S)$, the transition system of all *synchronous* executions
2. Check for safety on $TS(S)$ to
 - ▶ ensure equivalence between original system and the projections of the choreography,
 - ▶ guarantee safety (no deadlock, no orphan message)
3. Build a choreography (global graph) from $TS(S)$, relying on
 - ▶ the theory of regions, and
 - ▶ safe Petri nets.

1. Synchronous Transition System of CFSMs



Synchronous Transition System (TS(S))



2. Check for Safety: *Generalised Multiparty Compatibility (GMC)*

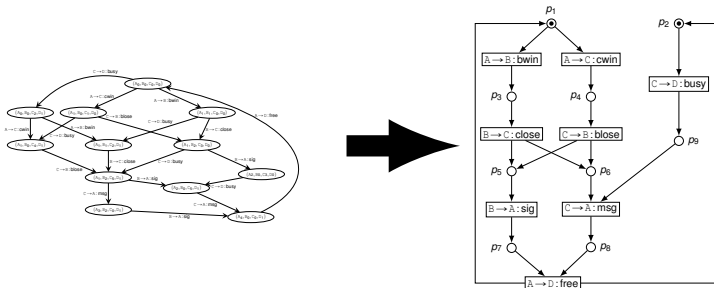
1. Representability: ensures that TS contains enough information
 - Essentially: the projection of TS onto a participant must be branching bisimilar to the original machine.
2. Branching Property: each branching in TS is “well-formed”
 - it either corresponds to concurrent interactions, or
 - each participant eventually knows which branch was chosen.



3. Build a global graph

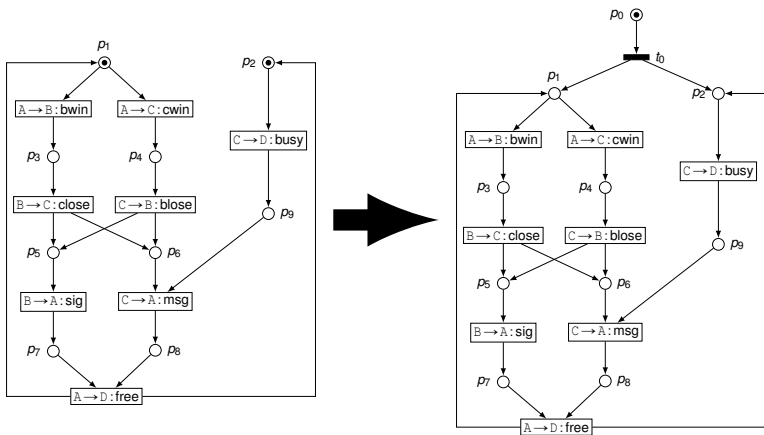


1: From TS to Petri Net

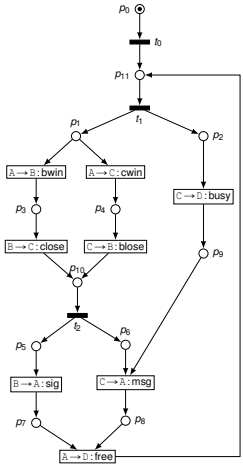
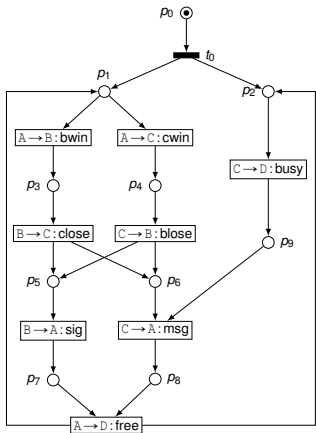


We use the work of Cortadella et al. (1998), based on the **theory of regions**, to derive a *safe* and *extended free-choice* Petri net from the Synchronous Transition System.

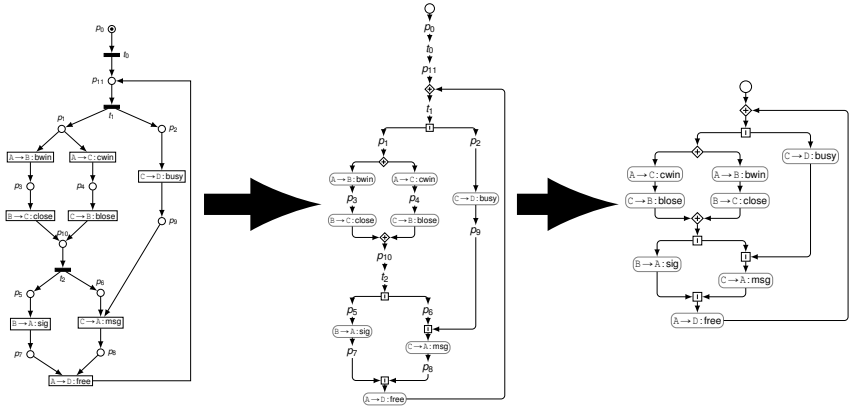
2: From Petri Net to One-source Petri Net



3: From One-source Petri Net to Joined Petri Net



4 & 5 : From Jointed Petri Net to Global Graph



Prototype Implementation

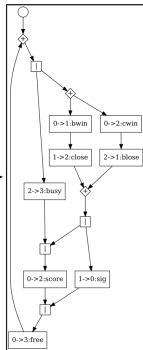
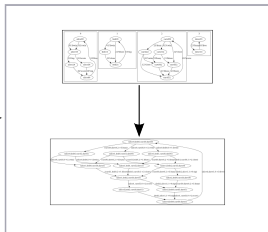


```
.outputs
.state graph
alice0 1 ! bwin alice1
alice0 2 ! cwin alice1
alice1 1 ? sig alice2
alice1 2 ! score alice2
alice2 1 ! score alice2
alice3 1 ? sig alice4
alice4 3 ! free alice0
.marking alice0
.end

.outputs
.state graph
bob0 0 ? bwin bob1
bob1 2 ! close bob2
bob0 2 ? blose bob2
bob2 0 ! sig bob0
.marking bob0
.end

.outputs
.state graph
carol0 1 ? close carol1
carol0 3 ! busy carol2
carol0 0 ? cwin carol1
carol1 3 ! busy carol4
carol1 1 ! blose carol3
carol2 1 ? close carol3
carol2 0 ? cwin carol4
carol3 3 ! busy carol5
carol4 1 ! blose carol5
carol5 0 ? score carol0
.marking carol0
.end

.outputs
.state graph
dave0 2 ? busy dave1
dave1 0 ? free dave0
.marking dave0
.end
```



<https://bitbucket.org/julien-lange/gmc-synthesis>



Conclusions and Future Work

- ▶ Summing up:
 - ▶ An effective way of checking properties of CFSMs, and whether one can construct a global graph from them.
 - ▶ An algorithm based on the theory of regions.
 - ▶ A CFSMs characterisation of well-formed *graphical global types*.
 - ▶ Tool(s):
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- ▶ Recent extension to support **Communicating Timed Automata**:
 - ▶ L. Bocchi, J. L., and N. Yoshida: Meeting Deadlines Together.
 - ▶ To appear in CONCUR'15
 - ▶ Preprint: <http://www.doc.ic.ac.uk/~jlange/cta/>



Thanks!

Any questions?

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